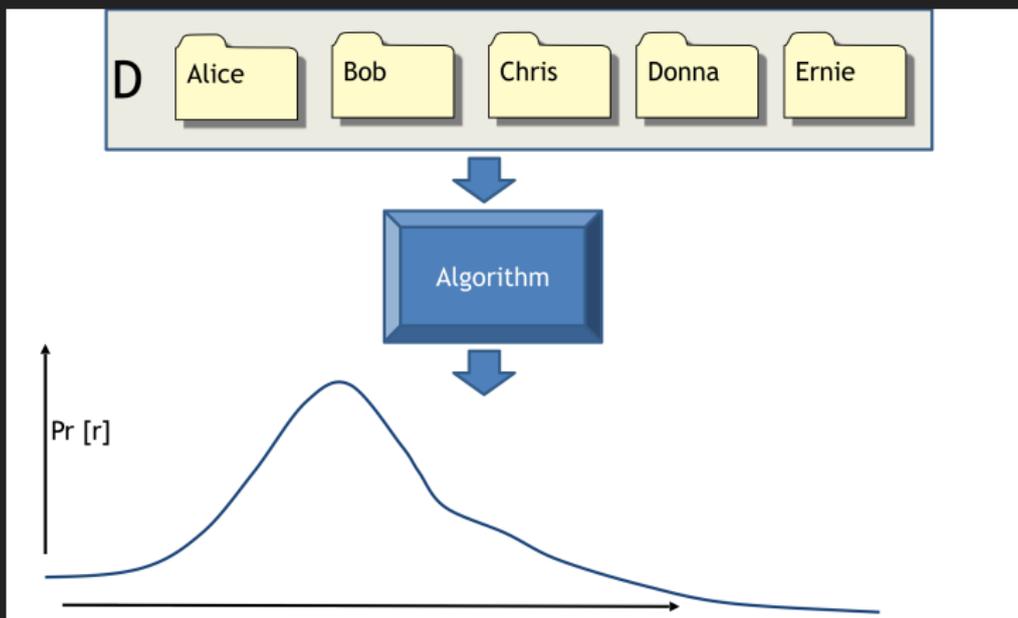


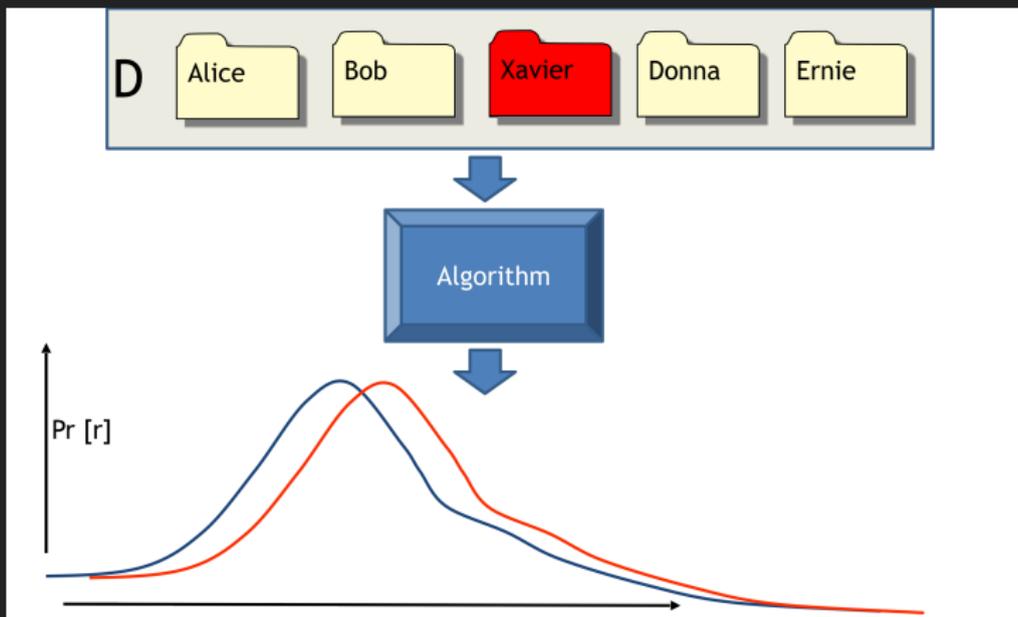
★-Liftings for Differential Privacy and f -Divergences

Gilles Barthe, Thomas Espitau,
Justin Hsu, Tetsuya Sato, Pierre-Yves Strub

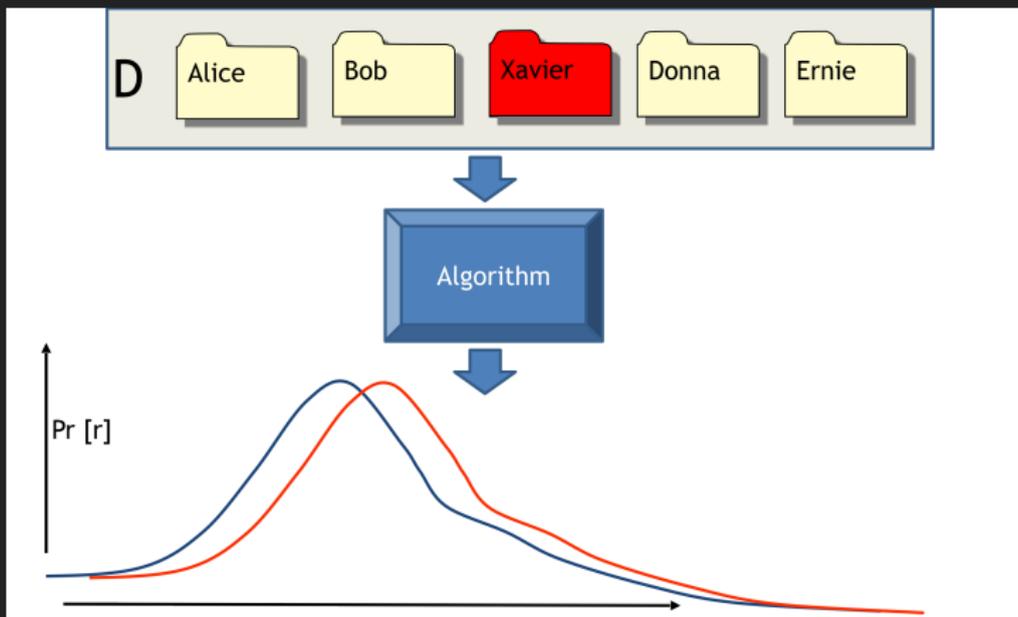
Differential privacy: probabilistic program property



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Output depends **only a little**
on any **single** individual's data

More formally

Definition (Dwork, McSherry, Nissim, Smith)

An algorithm is (ϵ, δ) -differentially private if, for every two adjacent inputs, the output distributions μ_1, μ_2 satisfy:

$$\Delta_\epsilon(\mu_1, \mu_2) \leq \delta \triangleq \text{for all sets } S, \mu_1(S) \leq e^\epsilon \cdot \mu_2(S) + \delta$$

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Behaves well under composition: “ ϵ and δ add up”

Sequentially composing an (ϵ, δ) -private program and an (ϵ', δ') -private program is $(\epsilon + \epsilon', \delta + \delta')$ -private.

How to verify this property?

Use ideas from probabilistic bisimulation

- ▶ $\Delta_\epsilon(\mu_1, \mu_2) \leq \delta$ means “approximately similar”
- ▶ Composition \iff approximate probabilistic bisimulation

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Foundation for many styles of program verification

- ▶ Linear and dependent type systems
- ▶ Product program constructions
- ▶ Relational program logics

Review: Probabilistic Liftings and Approximate Liftings

Probabilistic liftings

Lift a binary relation R on pairs $S \times T$
to a relation $\langle R \rangle$ on distributions $\text{Distr}(S) \times \text{Distr}(T)$

Definition (Larsen and Skou)

Let $R \subseteq S \times T$ be a relation. Two distributions are related $\mu_1 \langle R \rangle \mu_2$ if there exists a **witness** $\eta \in \text{Distr}(S \times T)$ such that:

1. $\pi_1(\eta) = \mu_1$ and $\pi_2(\eta) = \mu_2$,
2. $\eta(s, t) > 0$ only when $(s, t) \in R$.

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Example

$\mu_1 \langle = \rangle \mu_2$ is equivalent to $\mu_1 = \mu_2$.

An equivalent definition via Strassen's theorem

Theorem (Strassen 1965)

Let $R \subseteq S \times T$ be a relation. Then $\mu_1 \langle R \rangle \mu_2$ if and only if:

for all subsets $A \subseteq S$, $\mu_1(A) \leq \mu_2(R(A))$

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Approximate liftings

Intuition

- ▶ Approximately relate two distributions μ_1 and μ_2
- ▶ Add numeric indexes (ϵ, δ) to lifting

Want:

- ▶ Given $R \subseteq S \times T$, lift to $\langle R \rangle^{(\epsilon, \delta)} \subseteq \text{Distr}(S) \times \text{Distr}(T)$
- ▶ $\mu_1 \langle = \rangle^{(\epsilon, \delta)} \mu_2$ should be equivalent to $\Delta_\epsilon(\mu_1, \mu_2) \leq \delta$

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Previous definitions: “Existential”

Let $R \subseteq S \times T$ be a binary relation.
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One witness (Barthe, Köpf, Olmedo, Zanella-Béguelin)

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Two witnesses (Barthe and Olmedo)

There **exists** $\eta_L, \eta_R \in \text{Distr}(S \times T)$ such that

1. $\pi_1(\eta_L) = \mu_1$ and $\pi_2(\eta_R) = \mu_2$,
2. $\eta_L(s, t), \eta_R(s, t) > 0$ only when $(s, t) \in R$,
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No witnesses (Sato)

For all subsets $A \subseteq S$, we have

$$\mu_1(A) \leq e^\epsilon \cdot \mu_2(R(A)) + \delta$$

Which definition is the “right” one?

Definitions support different properties and constructions

	PW-Eq	Up-to-bad	Acc. Bd.	Subset	Mapping	Adv. Comp.
1-witness	?	?	Yes	?	?	?
2-witness	Yes	Almost*	No	Almost*	Almost*	Yes
Universal	Yes	Yes	Yes	Yes	Yes	?

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Broad tradeoff: How general?

- ▶ Less general: less compositional
- ▶ More general: harder to prove properties about

Our work: \star -Liftings, Equivalences,
and an approximate Strassen's theorem

New definition: \star -liftings

Generalize 2-witness lifting by adding a new point

Let $R \subseteq S \times T$ be a binary relation, and let $A^\star = A \cup \{\star\}$.
Two distributions are related by $\mu_1 \langle R^\star \rangle^{(\epsilon, \delta)} \mu_2$ if:

There **exists** $\eta_L, \eta_R \in \text{Distr}(S^\star \times T^\star)$ such that

1. $\pi_1(\eta_L) = \mu_1$ and $\pi_2(\eta_R) = \mu_2$,
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Intuition

- ▶ \star is a default point for tracking “unimportant” mass

Why is \star -lifting a good definition?

Previously known

One-witness $(??)$ Two-witness \Rightarrow Universal

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\star -liftings unify known approximate liftings

One-witness $\iff \star$ -lifting \iff Universal

Approximate version of Strassen's theorem

★-liftings are equivalent to “universal” approximate liftings

Theorem

Let S, T be discrete (countable) sets, and let $R \subseteq S \times T$ be a relation. Then $\mu_1 \langle R^* \rangle^{(\epsilon, \delta)} \mu_2$ if and only if:

for all sets $A \subseteq S$, $\mu_1(A) \leq e^\epsilon \cdot \mu_2(R(A)) + \delta$

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- ▶ Nodes

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- Outbound $c(\top, s)$ given by $\exp(-\epsilon) \cdot \mu_1$

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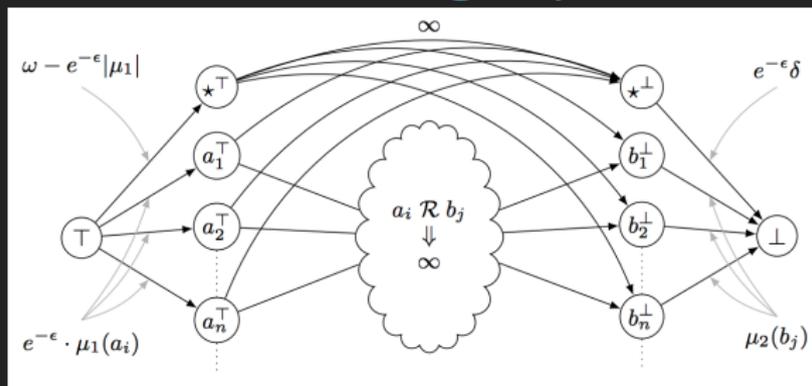
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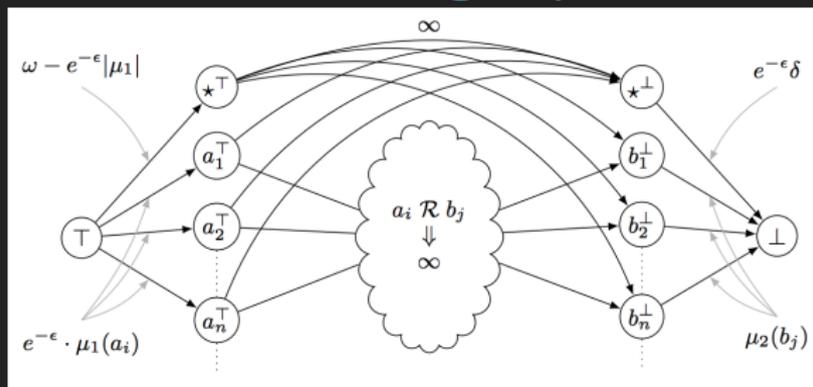
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- Outbound $c(\top, s)$ given by $\exp(-\epsilon) \cdot \mu_1$
- Incoming $c(t, \perp)$ given by μ_2

Proof sketch (universal lifting implies \star -lifting)



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Universal lifting \implies minimum cut large

- ▶ **Max-flow min-cut:** there is a **large** flow f from \top to \perp
- ▶ Use $f(s, t)$ to recover \star -lifting witnesses (η_L, η_R) , conclude:

$$\mu_1 \langle R^\star \rangle^{(\epsilon, \delta)} \mu_2$$

Other Results and Future Directions

See the paper for ...

- Further properties of \star -liftings
- Symmetric \star -liftings
and advanced composition
- \star -liftings for f -divergences

Wrapping up: Future directions and other speculation

Open questions

- ▶ Generalize to continuous distributions?
- ▶ Similar equivalences for other approximate lifting?
- ▶ Which properties should approximate liftings satisfy?

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Mild speculation

★-liftings are the “right” approximate version of probabilistic couplings

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